

FlexCap: Exploring Hardware Capabilities in Unikernels and Flexible Isolation Oses

CHERITech'24

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Objectives

- Investigate on Morello (real hardware), in the context of specialised minimal single address space operating systems (**unikernels**), the following:
 - Various approaches to **compartmentalisation in CHERI hybrid capability mode** and the resulting trade-offs in terms of performance, scalability, security, and engineering effort
 - The benefits of **safe C** obtained through the use of **CHERI pure capabilities**
 - How CHERI can help address one of the fundamental challenges of single address space operating systems: **multi-process applications support**

Outline

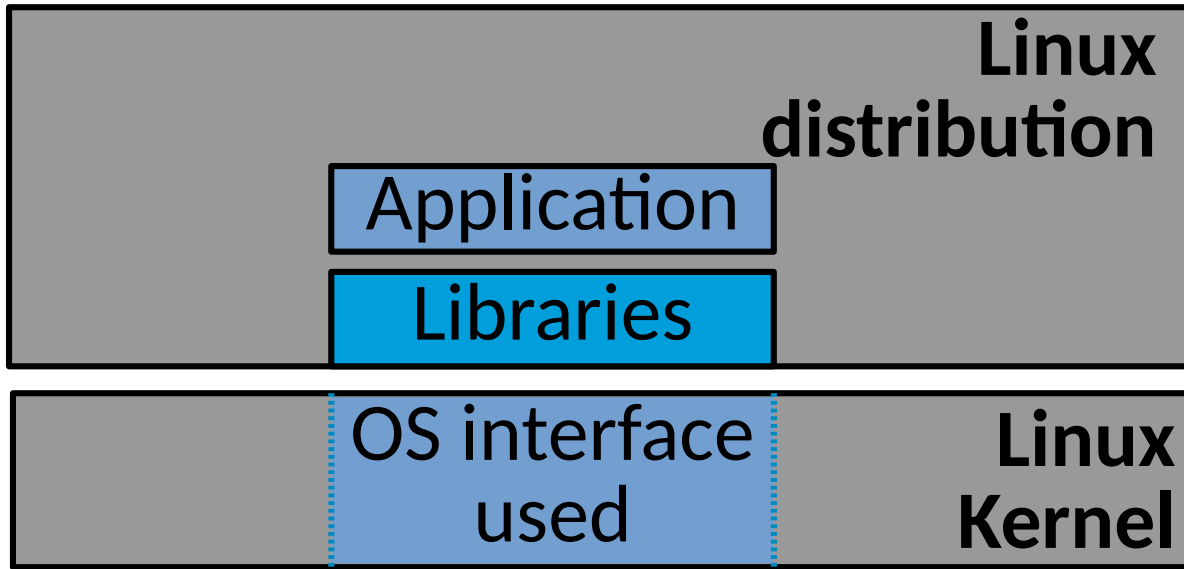
- 1) Unikernels
- 2) Progress on Unikernel Compartmentalisation
- 3) Progress on Purecap Unikernels
- 4) Progress on Support for Multi-Process Applications
- 5) Conclusion



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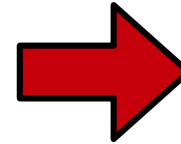
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Unikernels

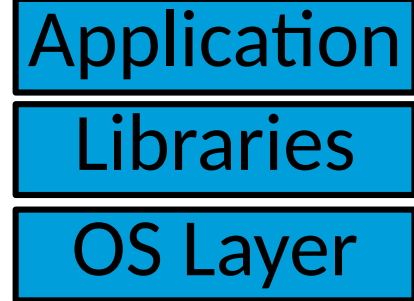
Full-fledged Virtual Machine



 Useful software
 Software bloat



Unikernel



Hypervisor

Hardware

Unikernels

- **Unikernel:** application + dependencies + thin OS layer compiled as a static binary running on top of a hypervisor
- **Single purpose OS:** 1 instance runs 1 application
 - OS can be specialised for the app (libOS/Exokernel model)
- **Lightweight:** fast boot time, low memory/disk footprint
 - Costs and attack surface reduction
- **Single binary and single address space** for the OS + application
 - *No isolation within the unikernel*, system calls are (fast) function calls



Unikernels + CHERI/Morello

- **Single address space** nature of unikernels aligns well with the protection model suggested by CHERI/Morello
- Today the **lack of isolation** inside a unikernel instance is concerning
- Motivated us to study bringing the security benefits of CHERI's compartmentalisation/safe C for unikernels while maintaining their lightweight and high-performance nature



+



arm

Morello Program

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Effort on Compartmentalisation

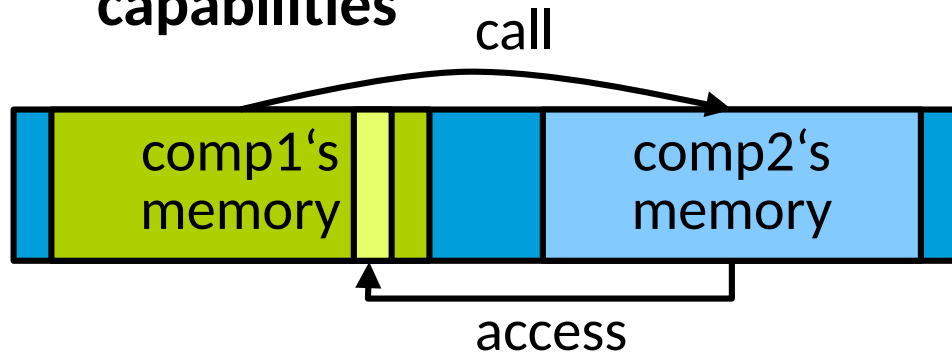
- Basic port of FlexOS¹ (compartmentalisation-aware version of the Unikraft² Unikernel) **to run bare metal on Morello A64**
- Development of compartmentalisation abstractions leveraging **hybrid mode** (for compatibility) with protection domains defined by DDC/PCC
- Development of **two methods of cross-compartment data sharing** trading off engineering effort/scalability to many compartments/security
- **Performance, security and engineering effort evaluation** of the prototype with 2 popular applications

¹Lefeuvre et al., “FlexOS: Towards Flexible OS Isolation”, ASPLOS’22

²Kuenzer et al. “Unikraft: fast, specialized unikernels the easy way”, EuroSys’21

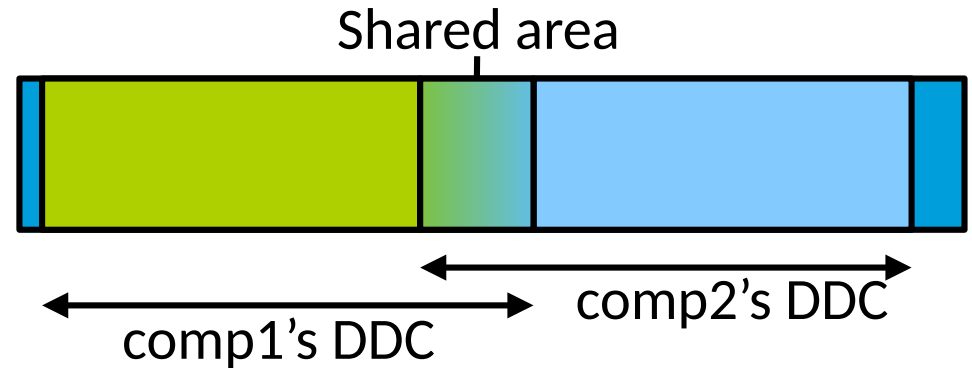
Compartmentalisation: Cross-Compartments Data Sharing

- Method 1: **pass shared data as capabilities**



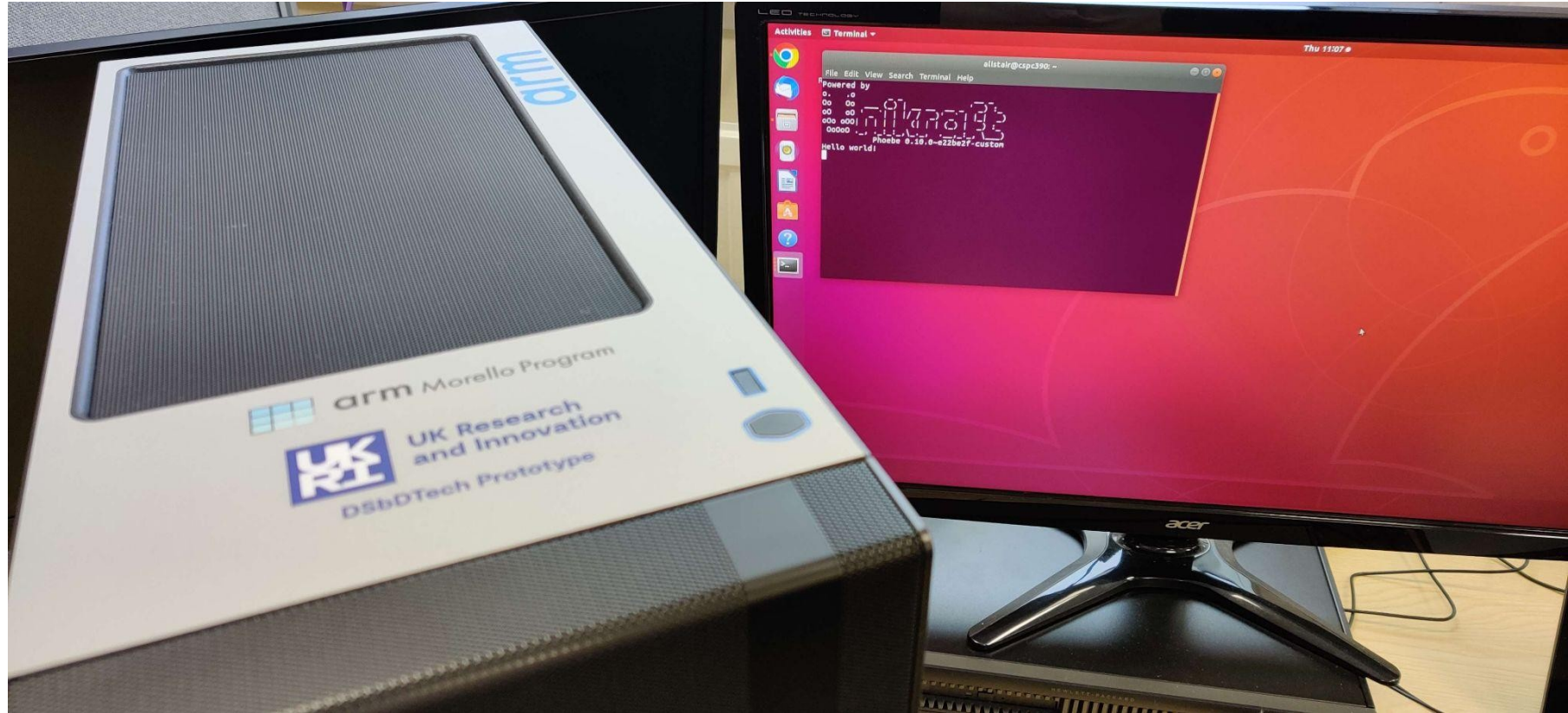
- Pros: security, scalability to high numbers of compartment
- Cons: **engineering effort**/scalability to large compartments

- Method 2: **overlapping DDCs**



- Pros: low engineering effort
- Cons: security, scalability to many compartments

Unikraft/FlexOS on Morello



<https://unikraft.org/blog/2022-12-01-unikraft-on-morello>

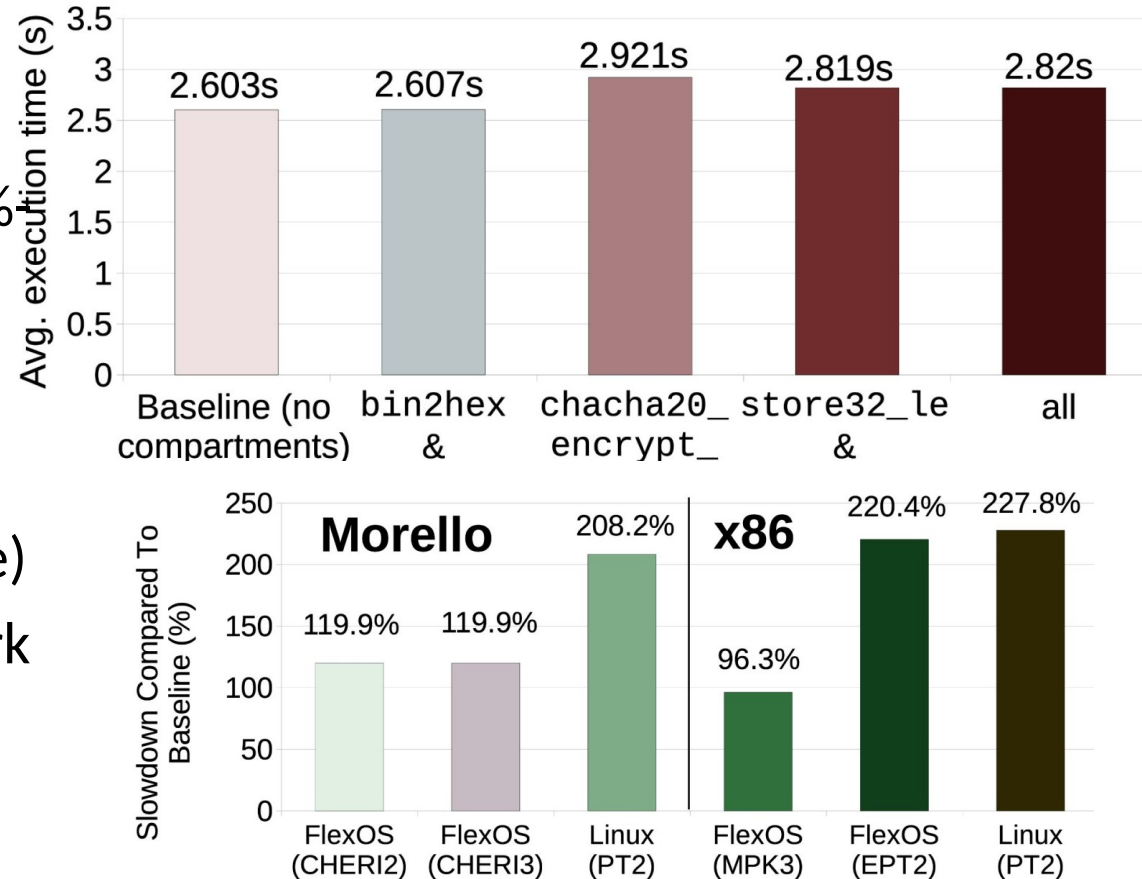
Compartmentalisation: Evaluation

Shared data as capabilities:

- With carefully selected functions overhead is low (Libsodium: 0.1% 12.2%)

Overlapping DDC:

- Performance overhead same order of magnitude to MPK and lower than EPT on FlexOS (SQLite)
- Runs faster than same benchmark on Linux with user/kernel isolation (SQLite)

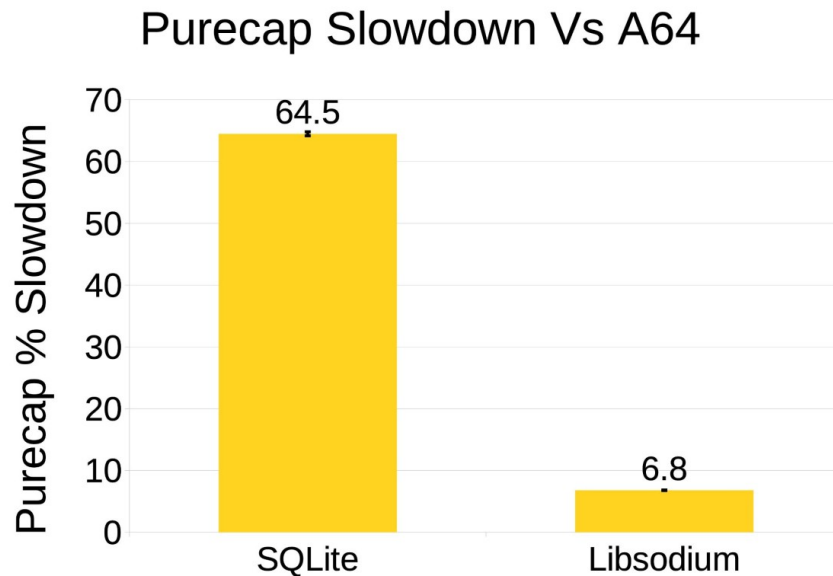


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Purecap Unikraft

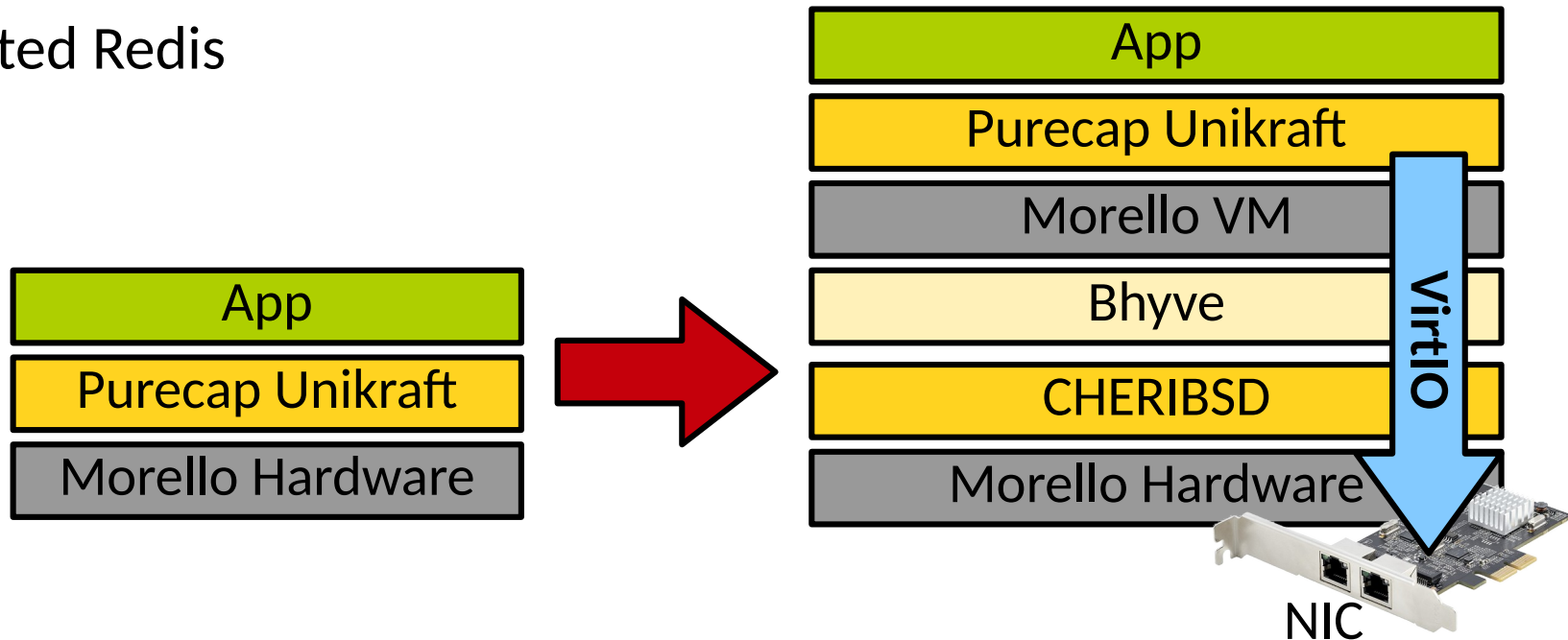
- We ported **Unikraft Unikernel OS** to run bare metal on Morello in **purecap mode**
 - Updates made to the platform code (boot process), memory allocator, and various other low-level subsystems (pointer arithmetics)
- We ported **libsodium and SQLite** to run on top of purecap Unikraft



| Benchmark | Avg. A64 cycles | Avg. Purecap cycles |
|-----------|-----------------|---------------------|
| SQLite | 2.6M | 4.2M |
| Libsodium | 130.2M | 139M |

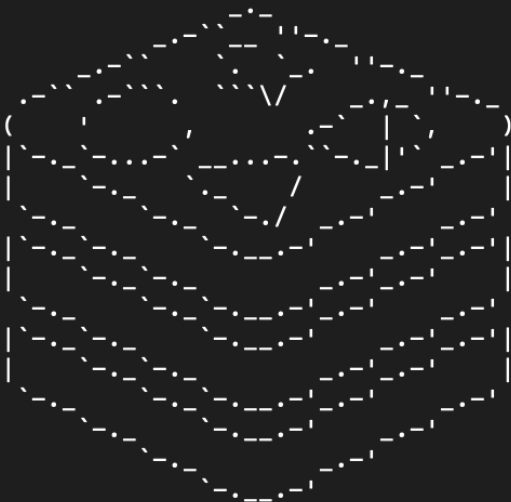
Beyond Bare Metal

- A custom OS on bare metal is limited by the lack **drivers for I/O**
- Ported Unikraft to run on top of the **bhyve hypervisor** and to use the **virtio-net paravirtualised network driver**
- Ported Redis



Beyond Bare Metal

```
CRIT: [libredis_server] Server config file /redis.conf
1:C 01 Jan 1970 00:00:00.091 # o000o000o000o Redis is starting o000o000o000o
1:C 01 Jan 1970 00:00:00.093 # Redis version=5.0.6, bits=64, commit=c5ee3442, modified=1, pid=1, just started
1:C 01 Jan 1970 00:00:00.094 # Configuration loaded
CRIT: [libredis_server] Pre init server
CRIT: [libredis_server] post init server
```



Redis 5.0.6 (c5ee3442/1) 64 bit

Running in standalone mode
Port: 6379
PID: 1

<http://redis.io>

```
1:M 01 Jan 1970 00:00:00.113 # Server initialized
1:M 01 Jan 1970 00:00:00.114 * Ready to accept connections
```


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SASOs & POSIX fork()

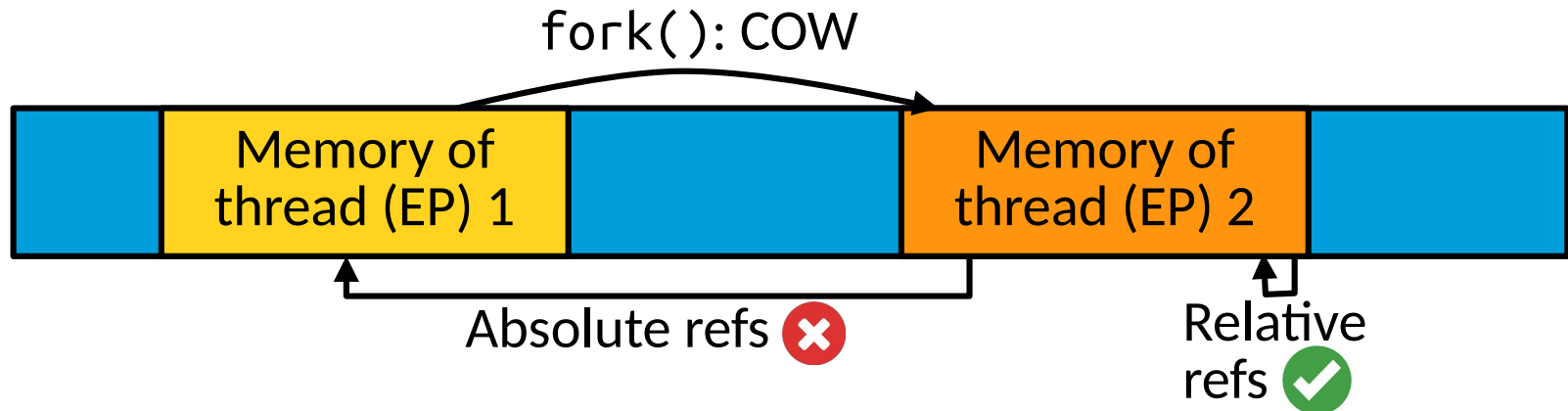
- Support for multi-process applications (i.e. POSIX fork()) is a well-know design limitation of single address space OSes (SASOs)
- Existing solutions for unikernels^{1,2} spawn 1 unikernel per process and implement IPCs in the hypervisor
 - This break the fundamental "single address space" nature of SASOs, loose some benefits
- **How can we support fork() within a single address space?**

¹Zhang et al., "KylinX: A Dynamic Library Operating System for Simplified and Efficient Cloud Virtualization", ATC'18

²Lupu et al., "Nephele: Extending Virtualization Environments for Cloning Unikernel-Based Vms", EuroSys'23

SASOs & POSIX fork()

- Key idea: **emulate processes with** threads
 - Locate the memory (code & data) relevant to each emulated process (EP) in a specific area
 - Upon fork, COW that space somewhere else and create another thread
 - **Challenges: inter-emulated process isolation, memory references**



SASOses & POSIX fork()


- **CHERI-powered solutions:**
 - Assume **PIE** to maximise relative memory references, and fixup absolute references on-demand during the COW by **scanning tagged memory to track pointers**
 - Inter-emulated processes **isolation based on purecap**, becomes a twofold problem:
 - 1) Ensure no capability leak between parent and child (i.e. properly fixup all absolute references during COW)
 - 2) Segregate and isolate the kernel's memory from the emulated processes as it is now an ambient authority
- Solutions in the process of being implemented

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Conclusion

- We look at various aspects of single address space OSes running on top of Morello:
 - Hybrid compartmentalisation
 - Safe C/purecap
 - Support for multiprocess applications
- Please come check out our poster on fork() support in SASOSes today
- And our PLOS'23 paper:
J. Kressel, H. Lefeuvre, P. Olivier, **Software Compartmentalization Trade-Offs with Hardware Capabilities**, PLOS'23
- <https://flexcap-project.github.io/>



Software Compartmentalization Trade-Offs with Hardware Capabilities

John Alistair Kressel, Hugo Lefeuvre, Pierre Olivier
The University of Manchester
Manchester, UK

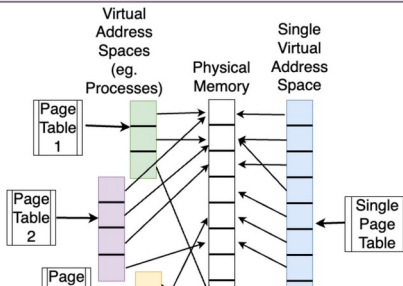
Abstract
Compartmentalization is a form of defensive software design in which an application is broken down into isolated but communicating components. This is achieved by enforcing the principle of least privilege [31]. Compartmentalization is a form of defensive software design in which an application is broken down into isolated but communicating components. This is achieved by enforcing the principle of least privilege [31].

1 Introduction
Software compartmentalization is one of the ways to enforce the principle of least privilege [31]. Compartmentalization is a form of defensive software design in which an application is broken down into isolated but communicating components. This is achieved by enforcing the principle of least privilege [31].

SFork: Supporting Complex Multi-Process Applications in a Single Address Space OS

John Alistair Kressel, Hugo Lefeuvre, Pierre Olivier
The University of Manchester

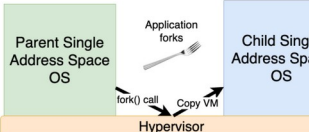
Single Address Space OSes (SASOS)



Problem: Lack of fork()

SASOSes are incompatible with forking, which creates a copy of their address space.

Most existing solutions treat OS as a single address space.



× Loses single address space performance during multiple address space creation.